On Routing Policies for Synchronized Queues

Mauro Escobar

Advisor: Mariana Olvera-Cravioto

Industrial Engineering and Operations Research Columbia University

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1 Motivation

2 Model

- 3 Preliminary Results
- 4 Numerical Results

5 Conclusions



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• Cloud computing

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- Large network of parallel servers

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- Large network of parallel servers
- Cost of information, single centralized queue is not scalable



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- All pieces must start service at the same time, FCFS

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Cost of Information vs. Service Quality?

We analize the **expected waiting time** in the 3 cases.















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Analyzed by Green [G80] with service time of pieces $S \sim \text{Exp}(\mu)$.







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Let f be a distribution of the number of pieces such that $f_n \stackrel{n}{\Rightarrow} f$, with finite mean $\mathsf{E}[N] = \sum_{k=1}^{\infty} kf(k)$.

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- Green [G80]: if $\lambda \mathsf{E}[N] / \mu < 1$, then $W^{(n)} \stackrel{n}{\Rightarrow} 0$.
- Olvera-Cravioto & Ruiz-Lacedelli [OR14]: $W^{(n)} \stackrel{n}{\Rightarrow} W$, if there exists $\beta > 0$ such that

$$\mathsf{E}\left[\sum_{i=1}^{N} e^{\beta(\chi_i - \tau_i)}\right] = \frac{\lambda \mathsf{E}\left[N\right]^2}{\beta + \lambda \mathsf{E}\left[N\right]} \mathsf{E}\left[e^{\beta\chi_1}\right] < 1,$$

where $N \sim f$, $\chi_i \sim S$ and $\tau_i \sim \text{Exp}(\lambda \mathsf{E}[N])$ are independent.



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- 30,000 jobs

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- Distribution of the number of pieces:
 - $N \sim \text{Poisson (light tail)}$
 - $N \sim$ Poisson composed with Pareto (heavy tail)
- Service time of the pieces:
 - $S \sim \text{Unif}(0, 1)$
 - $S \sim \operatorname{Exp}(1)$

Numerical Results: Example 1

Parameters:

$$\lambda = 0.1$$
 $S \sim \text{Unif}(0, 1)$ $\mathsf{E}[N] = 2 \text{ (light tail)}$

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Average Waiting Time:

 $8 \cdot 10^{-2}$ $6 \cdot 10^{-2}$ $4 \cdot 10^{-2}$ $2 \cdot 10^{-2}$ $\xrightarrow{} Opt \ \ \mathcal{C} Int$ $\xrightarrow{} Blind$ 0 0 200400 600 800 1,000 Number of Servers

 $\lambda = 0.1 \qquad S \sim \mathrm{Unif}(0,1) \qquad \mathsf{E}\left[N\right] = 20 \text{ (heavy tail)}$

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$$\lambda = 0.015$$
 $S \sim \text{Exp}(1)$ $\mathsf{E}[N] = 4 \text{ (light tail)}$

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 $\lambda = 0.015$ $S \sim \text{Exp}(1)$ $\mathsf{E}[N] = 4$ (light tail)

Average Waiting Time:



 $\lambda = 0.0025 \qquad S \sim \mathrm{Exp}(1) \qquad \mathsf{E}\left[N\right] = 10 \text{ (heavy tail)}$

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 $\lambda = 0.0025$ $S \sim \text{Exp}(1)$ $\mathsf{E}[N] = 10$ (heavy tail)

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- Light and heavy tails behave similarly
- Only knowing which servers are empty, decreases the waiting time significantly

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